

## Simple Example: Knapsack Problem

- We are given a set  $N = \{1, \dots, n\}$  of items and a capacity  $W$ .
- There is a profit  $p_i$  and a size  $w_i$  associated with each item  $i \in N$ .
- We want to choose the set of items that maximizes profit subject to the constraint that their total size does not exceed the capacity.
- The most straightforward formulation is to introduce a binary variable  $x_i$  associated with each item.
- $x_i$  takes value 1 if item  $i$  is chosen and 0 otherwise.
- Then the formulation is

$$\begin{aligned} & \max \sum_{j=1}^n p_j x_j \\ \text{s.t. } & \sum_{j=1}^n w_j x_j \leq W \\ & x_i \in \{0, 1\} \quad \forall i \end{aligned}$$

- Is this formulation correct?

## An Alternative Formulation

- Let us call a set  $C \subseteq N$  a *cover* if  $\sum_{i \in C} w_i > W$ .
- Further, a cover  $C$  is *minimal* if  $\sum_{i \in C \setminus \{j\}} w_i \leq W$  for all  $j \in C$ .
- Then we claim that the following is also a valid formulation of the original problem.

$$\begin{aligned} & \max \sum_{j=1}^n p_j x_j \\ \text{s.t. } & \sum_{j \in C} x_j \leq |C| - 1 \quad \text{for all minimal covers } C \\ & x_i \in \{0, 1\} \quad i \in N \end{aligned}$$

- Which formulation is “better”?